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Tardos's Fingerprinting Code over AWGN Channel

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Background

- Attack Model
- Hard & Soft Decision Method
- Estimation of Channel
- Proposed Tracing Algorithm
- Experimental Results
- Concluding Remarks

A server distributes personal copies of a content to N users.

c colluders mix their copies to forge a pirated copy.

The objective of a fingerprinting code

Identification of the colluders from the pirated copy.

c-secure code:

If the number of colluders is equal or less than c,

at least one of them can be identified.

Boneh Shaw (1995) concatenation code

Tardos (2003) Probabilistic fingerprinting code

Codeword of j-th user
$$\ m{x_j} = \{x_{j,1}, x_{j,2}, \dots, x_{j,L}\}$$

 $x_{j,i} \in \{0,1\}$ L: code length

When 4 users are colluded, they can produce a pirated copy under the following condition.

Codeword of
$$j$$
-th user $\ m{x_j} = \{x_{j,1}, x_{j,2}, \dots, x_{j,L}\}$
 $x_{j,i} \in \{0,1\}$ L: code length

When 4 users are colluded, they can produce a pirated copy under the following condition.

Each bit of the codeword is embedded into digital content assisted by a watermarking method.

Even if the watermarking method is robust, the extracted signal must be distorted by a noise.

Assumption additive white Gaussian noise (AWGN)



Т

Codeword of j-th user
$$x_j = \{x_{j,1}, x_{j,2}, \dots, x_{j,L}\}$$

 $x_{j,i} \in \{0,1\}$ L: code length

$$\begin{aligned} x_{j,i}, (1 \leq j \leq L) \text{ are independent.} \\ \Pr(x_{j,i} = 1) = p_i \quad \forall j \in [n] \end{aligned} \qquad \begin{array}{l} \text{i.i.d. random variables} \\ p_i \in [0, 1] \\ p_i \sim f(p) \end{aligned}$$

$$\begin{array}{ll} \textbf{racing Algorithm} \quad \boldsymbol{y} = \{y_1, y_2, \dots, y_L\} \text{ a pirated codeword} \\ \text{If the score } S_j \text{ exceeds a threshold } Z_i \text{ } j\text{-th user is judged guilty.} \\ \\ S_j = \sum_{i=1}^L y_i U_{j,i} \quad \text{where} \quad U_{j,i} = \begin{cases} \sqrt{\frac{1-p_i}{p_i}} & \text{if } x_{j,i} = 1 \\ -\sqrt{\frac{p_i}{1-p_i}} & \text{if } x_{j,i} = 0 \end{cases} \end{array}$$

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Symmetric version of the score

$$S_{j} = \sum_{i=1}^{L} y_{i} U_{j,i} \longrightarrow S_{j} = \sum_{i=1}^{L} (2y_{i} - 1) U_{j,i}$$

Gaussian approximation of the score's distribution based on CLT

innocent users
$$N(0,L)$$
 ϵ_1 : false-positive probability (Pr[FP])

Design of the threshold Z for a given ϵ_1

$$Z = \sqrt{2L} \cdot \operatorname{erfc}^{-1}\left(\frac{2\epsilon_1}{N}\right) \qquad N : \text{number of users}$$

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Correlation Score

Symmetric version of the score



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lacksimes The number $\,\widetilde{c}\,$ of colluders is not always equal and less than $\,c\,.$

The length L of codeword is fixed.

Our Goal Under a fixed length L, we want to catch as many colluders as possible without increasing the false-positive probability ϵ_1 no matter how many colluders are involved in.

Hard & Soft Decision Method



Hard & Soft Decision Method



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Comparison (SNR: 8dB)

code length :
$$L=10^4$$
 Pr[FP]: $\epsilon_1=10^{-4}$

 $\ensuremath{\texttt{\#}}$ of users : $N=10^4$



Although the HD method catches more colluders, the false-positive probability is increased.

Comparison ($\tilde{c} = 10$)

code length :
$$L=10^4$$
 Pr[FP]: $\epsilon_1=10^{-4}$

of users : ${\cal N}=10^4$



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Consideration

With the increase of noise, the false positive probability is increased.

Remark

In this experiment, the threshold Z_{HD} is calculated by a given ϵ_1 based on the CLT.

$$Z_{HD} = \sqrt{2L} \cdot \operatorname{erfc}^{-1}\left(\frac{2\epsilon_1}{N}\right)$$
$$= 3.97\sqrt{2L} \quad \text{(constant)}$$

 $Z \sim Z_{HD}$ (constant)

The false positive probability must be increased in the same fashion.

$$D \text{ is } 10^{-3} + 10^{-3} + 10^{-4} + 10^{$$

 10^{-1} c

False positive probability

$$\epsilon_1 = 10^{-4}, \ N = 10^4$$

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10 11

SNR [dB]

12

13

Proposed Method



Use only reliable elements for HD method (method I)



method I + unreliable elements (method II)

all users method I suspicious HD method with all elements users identified colluders

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Estimation of Channel

Estimating the variance
$$\sigma^2$$
 of noise $\hat{y} - \hat{y} - \hat{y} = \hat{y} + e^{N(0, \sigma^2)}$



Proposed Method I

Use only reliable elements for HD method (method I)



Proposed Method II



Proposed Method II



Comparison ($\tilde{c} = 10$)

code length : $L = 10^4$ Pr[FP] : $\epsilon_1 = 10^{-4}$

of users : $N = 10^4$

of bit flips :
$$P_{flip}L=1$$



The method II can detect many colluders without

increasing the false positive probability.

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The false positive probability is independent on the number of colluders.

- When the amount of noise is very small, the use of CLT is valid.
- The false positive probability is increased with the amount of noise. It is advisable to design a threshold considering the noise.
- We can identify colluders only from reliable elements without the serious increase of false positive probability. (method I)
- Among the suspicous users detected by the method I, the HD method can catch many colluders with less innocents.
- The false positive probability does not increase with the number of colluders for the method II.